According to the authors, in the Rate-Monotonic scheduling policy (RM policy), the task with the highest request rate is given higher priority. It is a fixed scheduling policy in that the priorities of the tasks are fixed. For RM scheduling policy, the *least upper bound* of processor utilization is given as , where is the number of tasks. Thus for RM policy, for more than one task, is less than one. For higher values of , it approaches . Also, for a given task set, if the processor utilization is greater than but less than one, then nothing can be said about the feasibility of the RM policy. This situation is overcome by using a dynamic scheduling policy, the *Deadline Driven Scheduling Algorithm*. In this, the tasks are assigned priorities dynamically based on their deadlines of their current requests. The task with the *earliest* deadline for its current request is assigned higher priority. For the deadline driven scheduling algorithm, also known as the *Earliest Deadline First* algorithm (EDF), the . Thus the algorithm assures feasibility of tasks if the processor utilization factor is less than or equal to one. Hence the EDF policy is globally optimum in that if a task set can be scheduled by any algorithm, then it can be scheduled by the EDF.